



HN-003-016304

Seat No. _____

M. Sc. (Mathematics) (Sem. III) (CBCS) Examination

May / June - 2017

MATH.CMT - 3004 : Discrete Mathematics

Faculty Code : 003

Subject Code : 016304

Time : $2\frac{1}{2}$ Hours]

[Total Marks : 70

- Instructions :** (1) Answer all the questions.
(2) Each question carries 14 marks.

1 Answer any Seven : 7×2=14

- (a) Verify that $00^*(0\vee 1)^*1$ is a regular expression over $\{0,1\}$.
- (b) Let A be a $m \times n$ Boolean matrix and B be a $n \times p$ Boolean matrix. Define the *Boolean product* of A and B .
- (c) Let R be the relation defined on $\mathbb{Z} \times \mathbb{Z}$ by $(a, b)R(x, y)$ if and only if $a - b = x - y$. Prove that R is an *equivalence relation* on $\mathbb{Z} \times \mathbb{Z}$.
- (d) If (P, \leq) is a partially ordered set, then prove that (P, \leq^{-1}) is a partially ordered set.
- (e) Define a *modular lattice* and illustrate it with an example.
- (f) Define a *Boolean Algebra* Let $n \geq 1$. Show that there exists a Boolean Algebra containing exactly 2^n elements.
- (g) Define *homomorphism of semigroups*. Let $(S, *) \rightarrow (T, *')$ be a surjective homomorphism of semigroups. If $(S, *)$ is a monoid, then show that $(T, *')$ is also a monoid.
- (h) Let $I = \{a, b\}$ and $L \subseteq I^*$, When is L said to be a *type three language* ?
- (i) Define a *machine congruence* on a finite state machine.
- (j) Define the *direct derivability* relation associated with a phrase structure grammar G .

- 2 Answer any Two :** **2×7=14**
- (a) Let H be a normal subgroup of a group G . Define the relation R on G by aRb if and only if $ab^{-1} \in H$. Prove that R is a congruence relation on G .
- (b) State and prove the *fundamental theorem of homomorphism of semi-groups*.
- (c) Let $n \geq 1$. Show that D_n , the lattice of positive divisors of n is distributive.

- 3** (a) Let A be a finite nonempty set. Define the concept of *regular expression over A* . **5**
- (b) Let (L_i, \leq_i) be lattices for each $i \in \{1, \dots, n\}$. Let $L = L_1 \times \dots \times L_n$. Prove that (L, \leq) is a lattice, where \leq is the product partial order on L . **5**
- (c) Let p, q be propositions. Construct a truth table for the statement $(p \wedge q) \vee \tilde{p}$. is the negation of P . **4**

OR

- 3** (a) Let R be a relation defined on a nonempty set A . Prove that R^∞ is the transitive closure of R . **5**
- (b) Let $f: L_1 \rightarrow L_2$ be a bijection, where (L_i, \leq_i) is a lattice for each $i \in \{1, 2\}$. Show that $f(a \vee b) = f(a) \vee f(b)$ for all $a, b \in L_1$ if and only if f and f^{-1} both preserve order. **5**
- (c) Construct the truth table for the Boolean function $f: B_3 \rightarrow B$ determined by the Boolean polynomial $p(x_1, x_2, x_3) = (x_1 \wedge x_2) \vee (x_1 \vee (x_2' \wedge x_3))$. **4**

- 4 Answer any Two :** **2×7=14**
- (a) Let (L, \leq) be a finite Boolean Algebra. Let $a \in L, a \neq 0$. Let $W = \{b \in L : b \text{ is an atom of } (L, \leq) \text{ with } b \leq a\}$. Prove that $a = \vee_{b \in W} b$.
- (b) State and prove the Pumping lemma.

- (c) Let M be a Moore machine with $S = \{s_0, s_1, s_2\}$, $I = \{a, b\}$. $f_a : S \rightarrow S$ equals the identity mapping on S . $f_b(s_0) = s_1$, $f_b(s_1) = s_2$, $f_b(s_2) = s_2$ and $T = \{s_2\}$. Find $L(M)$ and also determine a regular expression α over I such that $L(M)$ is the regular set corresponding to α .

5 Answer any Two : **2×7=14**

- (a) Let M be a Moore machine with S as its state set. Prove that the w -compatibility relation R defined on S is a machine congruence on M and $L(M) = L(M/R)$.
- (b) Let $n \geq 1$ Prove that any function $f : B_n \rightarrow B$ is produced by a Boolean expression.
- (c) Let (L, \leq) be a lattice. Show that (L, \leq) is distributive if and only if for all $a, b, c \in L$, $(a \wedge b) \vee (b \wedge c) \vee (c \wedge a) = (a \vee b) \wedge (b \vee c) \wedge (c \vee a)$.
- (d) Let M be a Moore machine with I as its input set. Show that there exists a type three phrase structure grammar G with I as its set of terminal symbols such that $L(M) = L(G)$.